

Vanishing of Littlewood-Richardson polynomials is in P

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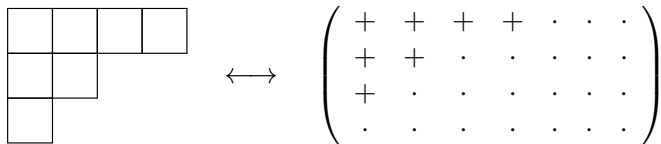
Background

Let $\lambda = (\lambda_1 \geq \lambda_2 \geq \cdots \geq \lambda_n \geq 0)$ be a partition with n nonnegative parts.

Definition

The **initial diagram** for λ is a grid with n rows and $m \geq n + \lambda_1 - 1$ columns in which the Young diagram of λ is in the northwest corner.

For example, $\lambda = (4, 2, 1, 0)$



Background

Definition

A **local move** is a mutation of any 2×2 subsquare of the form

$$\begin{array}{cc} + & \cdot \\ \cdot & \cdot \end{array} \mapsto \begin{array}{cc} \cdot & \cdot \\ \cdot & + \end{array}$$

The configurations of $+$'s in the grid resulting from local moves on the initial diagram for λ are called **plus diagrams**.

For example,

$$\left(\begin{array}{ccccccc} + & + & + & \cdot & \cdot & \cdot & \cdot \\ + & + & \cdot & \cdot & + & \cdot & \cdot \\ + & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \end{array} \right), \quad \left(\begin{array}{ccccccc} + & + & \cdot & \cdot & \cdot & \cdot & \cdot \\ + & \cdot & \cdot & + & \cdot & \cdot & \cdot \\ \cdot & \cdot & + & \cdot & \cdot & + & \cdot \\ \cdot & + & \cdot & \cdot & \cdot & \cdot & \cdot \end{array} \right)$$

Background

To a diagram T , we can assign weights,

$$wt_x(T) = \prod_i x_i^{+ \text{'s in row } i}$$

Or finer,

$$wt_{x,y}(T) = \prod_{(i,j) \in T} (x_i - y_j)$$

For example, for the following $T \in \text{Plus}((3, 2, 0))$,

$$\begin{pmatrix} + & \cdot & \cdot & \cdot & \cdot \\ + & \cdot & + & \cdot & \cdot \\ \cdot & \cdot & + & \cdot & + \end{pmatrix}$$

$$wt_x(T) = x_1 x_2^2 x_3^2$$

$$wt_{x,y}(T) = (x_1 - y_1)(x_2 - y_1)(x_2 - y_3)(x_3 - y_3)(x_3 - y_5).$$

Background

We can realize the Schur function as

$$s_\lambda(X) = \sum_{Plus(\lambda)} wt_x(T)$$

which forms a \mathbb{Z} -linear basis of $Sym[X]$ and similarly the factorial Schur function as

$$s_\lambda(X; Y) = \sum_{Plus(\lambda)} wt_{x,y}(T)$$

which forms a $\mathbb{Z}[Y]$ -linear basis of $Sym[X] \otimes_{\mathbb{Z}} \mathbb{Z}[Y]$.

Structure Coefficients

Then we can discuss their structure coefficients:

$$s_\lambda(X)s_\mu(X) = \sum_{\nu} c_{\lambda,\mu}^{\nu} s_{\nu}(X)$$

where $c_{\lambda,\mu}^{\nu} \geq 0$ is called the **Littlewood-Richardson coefficient**.
Extending this,

$$s_\lambda(X; Y)s_\mu(X; Y) = \sum_{\nu} C_{\lambda,\mu}^{\nu}(Y) s_{\nu}(X; Y)$$

where $C_{\lambda,\mu}^{\nu}$ is called the **Littlewood-Richardson polynomial**.
For example, using $\lambda = \mu = (1, 0)$ we can express

$$\begin{aligned} s_{(1,0)}(x_1, x_2; Y)^2 &= s_{(2,0)}(x_1, x_2; Y) + s_{(1,1)}(x_1, x_2; Y) \\ &\quad + (y_3 - y_2)s_{(1,0)}(x_1, x_2; Y). \end{aligned}$$

Complexity of $c_{\lambda,\mu}^\nu$

Theorem (DeLoera, McAllister, 2006; Mulmuley, Narayanan, Sohoni, 2012)

$c_{\lambda,\mu}^\nu \neq 0$ can be decided in polynomial time.

Theorem (Adve, Robichaux, Yong, 2017)

$c_{\lambda,\mu}^\nu(Y) \neq 0$ can be decided in polynomial time.

Eigenvalue Problem

Problem

For A, B, C $r \times r$ Hermitian matrices with real eigenvalues λ, μ, ν , how does $A + B = C$ constrain $(\lambda, \mu, \nu) \in (\mathbb{R}^r)^3$?

Horn (1962) conjectured a recursive list of inequalities, and Klyachko (1998) proved another list of inequalities. Klyachko's inequalities are satisfied by $(\lambda, \mu, \nu) \in \mathbb{Z}_{\geq 0}^{3r} \iff c_{N\lambda, N\mu}^{N\nu} \neq 0$ for some N , which imply those of Horn given another result.

Theorem (Knutson-Tao, 1999)

$c_{\lambda, \mu}^{\nu} \neq 0 \iff c_{N\lambda, N\mu}^{N\nu} \neq 0$ for all $N \in \mathbb{Z}_{>0}$.

These inequalities precisely control the vanishing of $c_{\lambda, \mu}^{\nu}$.

A Variation on the Eigenvalue Problem

Problem

For A, B, C Hermitian matrices with real eigenvalues λ, μ, ν , how does $A + B \geq C$ constrain $(\lambda, \mu, \nu) \in (\mathbb{R}^r)^3$?

Friedland (2000) proved inequalities that constrain those eigenvalues.

Theorem (Anderson-Richmond-Yong, 2013)

$C_{\lambda, \mu}^{\nu} \neq 0 \iff C_{N\lambda, N\mu}^{N\nu} \neq 0$ for all $N \in \mathbb{Z}_{>0}$.

Combining this result with Friedland's, they show that Friedland's inequalities control the vanishing of $C_{\lambda, \mu}^{\nu}$.

Proving the Complexity of $C_{\lambda,\mu}^\nu$

Our proof is a modification of the argument of Mulmuley, Narayanan, Sohoni.

We use a combinatorial rule for $C_{\lambda,\mu}^\nu$ in terms of edge-labeled tableaux of Thomas and Yong to construct a polytope $P_{\lambda,\mu}^\nu$ where

- $P_{\lambda,\mu}^\nu \cap \mathbb{Z}^{2\ell(\nu)\ell(\mu)} \neq \emptyset$ if and only if $C_{\lambda,\mu}^\nu > 0$, and
- $NP_{\lambda,\mu}^\nu = P_{N\lambda,N\mu}^{N\nu}$.

For example, T is an edge-labeled tableau that would be detected by $C_{(4,4,3),(4,2,1)}^{(3,2,2)}$ is

$$T = \begin{array}{|c|c|c|c|} \hline & & & \\ \hline & & 1 & 1 \\ \hline & 1 & 2 & \\ \hline & 2 & 3 & \\ \hline & 3 & & \\ \hline \end{array}$$

Proof Sketch

- $P_{\lambda,\mu}^\nu \neq \emptyset \iff C_{\lambda,\mu}^\nu \neq 0$.
(\Rightarrow) $P_{\lambda,\mu}^\nu \neq \emptyset$ implies there exists a rational vertex.
Then for some N , $NP_{\lambda,\mu}^\nu = P_{N\lambda,N\mu}^{N\nu}$ has a lattice point.
By construction of $P_{N\lambda,N\mu}^{N\nu}$, $C_{N\lambda,N\mu}^{N\nu} \neq 0$.
By equivariant saturation, $C_{\lambda,\mu}^\nu \neq 0$.
- Considering the linear programming problem and any objective function, we need to know the feasibility.
 - Simplex method has worst case exponential complexity.
 - Ellipsoid method has polynomial time complexity.

Conclusion and Summary

- The factorial schur functions $s_\lambda(X; Y)$ enrich the ordinary schur functions $s_\lambda(X)$.
- Through a variation on the eigenvalue problem, we discussed a connection between eigenvalues λ, μ, ν of matrices $A + B \geq C$ and when $C_{\lambda, \mu}^\nu \neq 0$.
- We proved $C_{\lambda, \mu}^\nu \neq 0$ is polynomial time to decide by combining combinatorial rule with equivariant saturation and integer programming.